# CS 357 D

Lecture 8

Orders and Lattices

http://cs357d.stanford.edu/

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# LATTICES AND ORDER

Supplementary Notes based on

Introduction to Lattices and Order by B.A. Davey and H.A. Priestley

Cambridge University Press, 2001

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#### Order

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Let P be a set. An order (or partial order) on P is a binary relation  $\leq$  on P such that for all  $x,y,z \in P$ :

- (i)  $x \le x$  (reflexivity)
- (ii)  $x \le y$  and  $y \le z$  implies  $x \le z$  (transitivity)
- (iii)  $x \le y$  and  $y \le x$  implies x = y (antisymmetry)

The relation  $\leq$  gives rise to the relation < of strict inequality:

$$x < y$$
 in P iff  $x \le y$  and  $x \ne y$ 

### Partially ordered set (Poset)

A set P equipped with an order relation  $\leq$  is called a partially ordered set, or poset

```
Example: (P, \leq) with P = \{ \perp, neg, zero, pos, T \} \leq = \{ (\bot, \bot), (neg, neg), (zero, zero), (pos, pos), (T, T), (\bot, neg), (\bot, zero), (\bot, pos), (\bot, T), (neg, T), (zero, T), (pos, T) \}
```

note that the elements neg , zero , and pos are not related to each other: we don't have neg  $\leq$  zero nor zero  $\leq$  neg

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#### Covering relation

 $(P, \leq)$ : ordered set  $x,y \in P$ 

x is covered by y , written x < y if

(i) x < y, and

intuitive meaning: there is

(ii)  $x \le z < y$  implies z = x

nothing between x and y

For a finite set, the ordering relation is determined by the covering relation and v.v.

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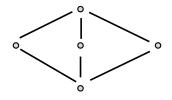
# Hasse diagrams

Hasse diagrams are a pictorial representation of the covering relation:

if x ⊰ y

x and y are connected by an edge, and x is drawn below y

Example:



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### Covering relation

Examples:

$$(N, \leq)$$
  $x \leq y$  if  $y = x + 1$ 

 $(\Re, \leq)$  no covering relation

( 
$$\wp(X)$$
 ,  $\subseteq$  ) A  $\prec$  B if B = A  $\cup$  {b} for some b  $\in$  X / A powerset of X

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0

### Special partially ordered sets

#### Chain

An ordered set P is a chain if for all  $x, y \in P$  either  $x \le y$  or  $y \le x$  (all elements are comparable)

Also known as: ▶ linearly ordered set 
▶ totally ordered set

Examples:

( Z ,  $\leq$  ) (set of all integers with the standard order) (  $\{\bot, \top\}$  ,  $\{(\bot, \bot), (\top, \top)\}$  )

#### Special partially ordered sets

#### **Antichain**

An ordered set P is an antichain if for all  $x, y \in P$  if  $x \le y$  then x = y (no elements are comparable)

0 0 0 0 0 0 0

### Examples:

```
( Z, {(x,x) | x ∈ Z}) (set of all integers with reflexive relation)
( { neg , zero , pos },
 {(neg , neg), (zero , zero), (pos , pos)})
```

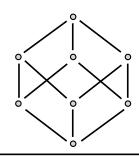
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## Hasse diagrams -- Example

Diagram of  $( \wp(\{1,2,3\}),\subseteq):$ 



3-element set

2-element sets

1-element sets

0-element set

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# Maps between orders

( P ,  $\leq_P$  ) , ( Q ,  $\leq_Q$  ) : ordered sets

 $f: P \rightarrow Q$  , function from P to Q , is

(i) monotone (order-preserving) if

$$x \leq_P y$$
 implies  $f(x) \leq_Q f(y)$ 

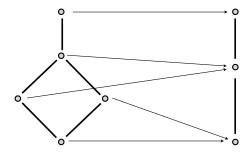
(ii) an order-embedding if

$$x \leq_P y$$
 iff  $f(x) \leq_Q f(y)$ 

(iii) an order-isomorphism if

$$x \leq_P y$$
 iff  $f(x) \leq_Q f(y)$  and f is onto

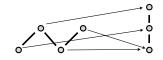
### Maps between orders -- example



monotone, but not an order-embedding

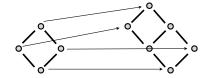
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### Maps between orders -- example



monotone

not order-embedding



monotone order-embedding

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### Top and Bottom

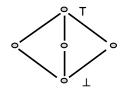
 $(P, \leq_P)$ : ordered set  $x \in P$ 

x is a bottom (least) element ( $\bot$ ) of P if  $\forall y \in P . x \leq_P y$ 

x is a top (greatest) element (T) of P if  $\forall y \in P . y \leq_P x$ 

top and bottom may not exist

top and bottom are unique if they exists



no top, no bottom

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### Top and bottom -- examples

( ℘ ( X ) , ⊆ ):

 $\perp = \emptyset$ 

T = X

 $(\{n \mid n \ge 0\}, \le) : \bot = 0$ 

no top

 $( \{ x \in R \mid a \le x \le b \}, \le ) : \bot = a$ 

 $(\{x \in R \mid a < x < b\}, \leq) :$  no bottom

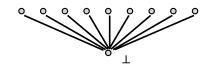
no top

### Lifting

Add a bottom element to an otherwise unordered set

S





 $S_{\perp} = S \cup \{ \perp \}$ 

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#### Maximal and minimal elements

 $(P, \leq_P)$ : ordered set,  $Q \subseteq P$ ,  $x \in Q$ 

x is a maximal element of Q if

for all  $y \in Q$ :  $x \le y$  implies x = y

x is a minimal element of Q if

2 maximal elements

for all  $y \in Q$ :  $y \le x$  implies x = y

Example:

(  $\wp(N)\N$ ,  $\subseteq$  ) has maximal elements  $N\n$  for all  $n \in N$ 

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# Upper bound

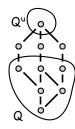
 $(P, \leq_P)$ : ordered set,  $Q \subseteq P, x \in P$ 

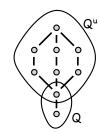
x is an upper bound of Q if  $\forall y \in Q . y \leq_P x$ 

 $Q^u: \{ x \in P \mid \forall y \in Q : y \leq_P x \}$  all upper bounds of Q

if Qu has a least element x:

x is called the least upper bound (lub) or supremum (sup) of Q





if Q has a top element T: sup Q = T

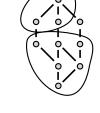
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#### Down-sets and Up-sets

( P ,  $\leq_P$  ): ordered set, Q  $\subseteq$  P

Q is an up-set (order filter, increasing set) if for all  $x \in Q$ ,  $y \in P$ :  $x \le y$  implies  $y \in Q$  (Q is closed under going up)

Q is a down-set ( order ideal , decreasing set ) if for all  $x \in Q$ ,  $y \in P$ :  $y \le x$  implies  $y \in Q$  (Q is closed under going down)



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#### Lower bound

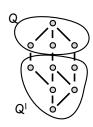
( P ,  $\leq_P$  ): ordered set, Q  $\subseteq$  P , x  $\in$  P

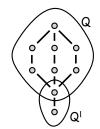
x is a lower bound of Q if  $\forall y \in Q . x \leq_P y$ 

 $Q^{l}: \{ x \in P \mid \forall y \in Q : x \leq_{P} y \}$  all lower bounds of Q

if Q1 has a greatest element x:

x is called the greatest lower bound (glb) or infimum (inf) of Q





if Q has a bottom element  $\bot$ : inf Q =  $\bot$ 

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### Upper and lower bounds

( P ,  $\leq_P$  ) : ordered set, assume P has  $\bot$  and T

$$\sup P = T$$
 inf  $P = \bot$ 

$$\sup \emptyset = \bot$$
 inf  $\emptyset = T$  Note:  $\emptyset^u = \emptyset^l = P$ 

#### Notation:

$$\sup(\{x,y\}) = x \lor y$$
 join  $\sup Q = \bigvee Q$   
 $\inf(\{x,y\}) = x \land y$  meet  $\inf Q = \bigwedge Q$ 

### Some properties:

if 
$$x \leq_P y$$
 then  $x \wedge y = x$  and  $x \vee y = y$ 

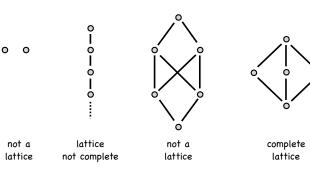
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#### Lattice

(P,  $\leq_P$ ): ordered set

P is a lattice if  $x \wedge y$  and  $x \vee y$  exist for all  $x, y \in P$ 

P is a complete lattice if  $\bigvee_Q$  and  $\bigwedge_Q$  exist for all  $Q\subseteq P$ 



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### Lattice

#### Note:

$$\varnothing\subseteq P$$
 ,  $P\subseteq P$  , so in a complete lattice 
$$\bigvee\varnothing$$
 ,  $\bigwedge\varnothing$  ,  $\bigvee P$  and  $\bigwedge P$  must all exist

recomplete lattice must be bounded

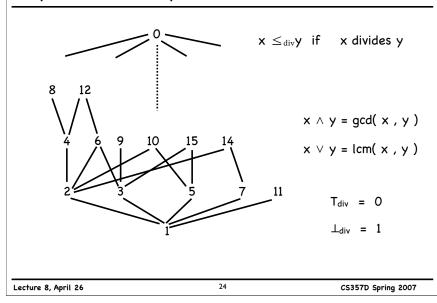
### Examples:

( 
$$\wp$$
 (  $X$  ) ,  $\subseteq$  ) is a complete lattice for any set  $X$ 

( N , 
$$\leq_{\text{div}}$$
 ) is a complete lattice divides relation

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### Complete lattice -- example



### Lattice as algebraic structure

 $(P, \leq_P)$ : lattice

Define  $\wedge$  and  $\vee$  as functions:  $\wedge$ ,  $\vee$ :  $P^2 \rightarrow P$ 

 $x \wedge y = \inf\{x, y\}$ 

 $x \lor y = \sup\{x, y\}$ 

Properties: ∧ ,∨ are order preserving

 $x \leq_{P} z$  implies  $x \wedge y \leq_{P} z \wedge y$  and  $x \vee y \leq_{P} z \vee y$ 

 $x \leq_P z$  and  $y \leq_P t$  implies  $x \wedge y \leq_P z \wedge t$  and

 $x \vee y \leq_{P} z \vee \uparrow$ 

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### Lattice as algebraic structure

Properties:

associative: 
$$(x \lor y) \lor z = x \lor (y \lor z)$$
  $(x \land y) \land z = x \land (y \land z)$ 

$$(x \wedge y) \wedge z = x \wedge (y \wedge z)$$

commutative:

$$x \wedge y = x \wedge y$$

$$x \lor y = x \lor y$$

idempotent:  $x \lor x = x$   $x \land x = x$ 

$$x \lor x = y$$

$$x \wedge x = x$$

absorption

$$x \land (x \lor y) = x$$
  $x \lor (x \land y) = x$ 

$$x \lor (x \land y) = x$$

In a lattice, join and meet are determined by the order and v.v.

$$(P, V_P, \Lambda_P)$$

 $(P, \leq_P)$   $(P, \vee_P, \wedge_P)$  can be used interchangeably

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# Lattice as algebraic structure

# Properties:

$$\bot \land a = \bot$$
  $\bot \lor a = a$ 

⊥ acts like 0

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 $T \wedge a = a$   $T \vee a = T$ 

$$T \vee a = 7$$

T acts like 1

Alternative representations:  $(P, \leq_P)$   $(P, \wedge, \vee)$ 

$$(P.<_{P})$$

Note:

in (N, lcm, gcd): 
$$0 = 1$$
 and  $1 = 0$ 

### Homomorphism

$$(P, \leq_P)$$
,  $(Q, \leq_Q)$ : ordered sets

$$f: P \rightarrow Q$$
 is a lattice homomorphism if

$$f\left( \ x \ \lor_P \ y \ \right) \ = \ f(\ x \ ) \ \lor_Q \ f(\ y \ ) \qquad \text{join preserving}$$

$$f(x \wedge_P y) = f(x) \wedge_Q f(y)$$
 meet preserving

f is a lattice homomorphism iff it is an order isomorphism

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#### Knaster-Tarski fixed point theorem

(P,  $\leq_P$ ): complete lattice

 $f: P \rightarrow P$  order preserving function (monotone)

greatest fixed point of f:  $gfp(f) = \bigwedge \{ x \in P \mid x \leq_P f(x) \}$ 

least fixed point of f:  $lfp(f) = \bigwedge \{ x \in P \mid f(x) \leq_P x \}$ 

if f is increasing ( $x \le_P f(x)$ ) lfp(f) can be obtained by starting with  $\bot$  and repeatedly applying f:

$$\bot_P \; \longrightarrow \; f(\bot_P) \; \longrightarrow \; f(f(\bot_P)) \; \longrightarrow \; f(f(f(\bot_P))) \; \longrightarrow \; \cdots$$

until a fixed point is reached:  $f^n(\bot_P) \, = \, f^{n+1}(\bot_P)$ 

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### Knaster-Tarski fixed point theorem

( $P, \leq_P$ ): complete lattice

 $f: P \rightarrow P$  order preserving function (monotone)

greatest fixed point of f:  $gfp(f) = \bigwedge \{ x \in P \mid x \leq_P f(x) \}$ 

least fixed point of f:  $lfp(f) = \bigwedge \{ x \in P \mid f(x) \leq_P x \}$ 

if f is decreasing  $(f(x) \le_P x)$  gfp(f) can be obtained by starting with T and repeatedly applying f:

$$\top_{P} \longrightarrow f(\top_{P}) \longrightarrow f(f(\top_{P})) \longrightarrow f(f(f(\top_{P}))) \longrightarrow \cdots$$

until a fixed point is reached:  $f^n(T_P) = f^{n+1}(T_P)$ 

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## Ascending and descending chain condition

( $P, \leq_P$ ): ordered set

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- (i) P has length n if the longest chain in P has length n
- (ii) P has finite length if all its chains are finite
- (iii) P satisfies the ascending chain condition (ACC) if for any sequence  $x_1 \leq_P x_2 \leq_P \dots$  of elements in P there exists k such that  $x_k = x_{k+1} = \dots$ .
- (iv) P satisfies the descending chain condition (DCC) if for any sequence  $x_1 \ge_P x_2 \ge_P \dots$  of elements in P there exists k such that  $x_k = x_{k+1} = \dots$

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Note: A finite lattice satisfies both the ACC and the DCC

# Ascending and descending chain condition: examples



length 3 satisfies both ACC and DCC

for finite set X of size n:  $( ( (X ), \subseteq ))$ length n+1
satisfies both ACC and DCC

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